

Eexam

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Advanced Computer Networking

Exam: IN2097 / Retake-Online Date: Monday 11th April, 2022

Examiner: Prof. Dr.-Ing. Georg Carle **Time:** 14:15 – 15:30

Working instructions

- This exam consists of 12 pages with a total of 5 problems.
 Please make sure now that you received a complete copy of the exam.
- The total amount of achievable credits in this exam is 75 credits.
- · Detaching pages from the exam is prohibited.
- · Allowed resources:
 - one analog dictionary English ↔ native language without annotations
 - the **provided cheatsheet** without annotations (print or digital copy)
- Subproblems marked by * can be solved without results of previous subproblems.
- Answers are only accepted if the solution approach is documented. Give a reason for each answer unless explicitly stated otherwise in the respective subproblem.
- Do not write with red or green colors nor use pencils.
- · Code of conduct:
 - I participate without the help of others and only use the allowed resources.
 - I do not share, discuss, or exchange any information related to the exam with anybody.
 - I feel in good health and I am able to participate in the exam.
 - I understood the examination policy, agree to the video supervision, and adhere to this process.

Problem 1 Quiz (15 credits)

The following questions cover multiple topics and can be solved independently of each other. The multiple choice questions need to be filled out as follows:

	Mark correct answers with a cr To undo a cross, completely fill	out the answer option	
	To re-mark an option, use a hu	-	
Each question has one s	ingle correct answer that is we	orth 1 credit.	
a)* What is the length of a	an IPv6 address?		
2 ¹²⁸ byte	☐ 128 byte	2 ¹²⁸ bit	16 byte
b)* Which of the following	IPv6 addresses is a Solicited N	Node Address?	
ff02::1:ffb9:fd7f	ff01::1:ffb9:fd7f	2002::1:bbb9:fd7f	ff02::1:bbb9:fd7f
, -	nly take to scan the complete IProcomplete address space is pro	•	and a rate of 10 000 packets
☐ 50 h	200 h	☐ 24 h	☐ 120 h
☐ In SDNs, a single co ☐ SDNs rely on a phys ☐ In SDNs, control pla ☐ An SDN is not requi	at a Software-Defined Network ontrol plane may control several sically centralized control plane ane and data plane are separate red to realize Network Function of part of the core P4 language	forwarding planes. e. Virtualization (NFV).	ation?
exact	range	ternary	Ipm
() + M/L'ala a a a a a l'a a a a a a	orlanda diban in language		
TCP BBR	rol algorithm is loss-based ?	☐ TCP Illinois	☐ TCP Reno
g)* What is the Organisat	ion Unique Identifier (OUI) and	at which position can it be fou	ind?
h)* Does the Spanning Tr	ee Protocol (STP) prevent routi	ng loops?	

i)* What is the main difference between classful and classless IP addressing, regarding prefix length?	
j)* Name the two latest DNS encryption mechanisms discussed in the lecture.	
k)* Explain where DNS encryption protocols introduced in the lecture are operating and why that is useful.	
I)* What is the difference between an active and an inactive timeout and when would they be triggered?	
m)* Why do QUIC ZMap scans require more bandwidth than TCP ZMap scans?	

Problem 2 Longest Prefix Matching (LPM) (14 credits)

Longest Prefix Matching (LPM) is the algorithm that is performed by a router. The router performs the LPM for each packet to determine the next hop of a packet and perform an appropriate forwarding decision.

There are serveral notations to represent the network parts of IPv4 addresses. The slash-notation /x, the dotted-decimal notation a.b.c.d, or a subnetmask represented as a hex value 0xfffffffff.

	a)* Convert /17 into the dotted-decimal notation.
l	b)* Convert 255.255.254.0 into a subnetmask as a hex value.
	c)* Explain briefly if 128.255.255.255 is a valid subnet mask.
1	A major part of the LPM is the matching operation (cf. Listing 1). The given matching function has three unsigned 32-bit integers as input arguments. The first argument contains the IPv4 address of the packet that should be forwarded (uint32 addr). The second argument contains the IPv4 address of a routing entry that should be matched (uint32 entry). The third argument contains the subnet mask of the routing entry (uint32 mask). match() returns true if the entry matches, false otherwise.
	boolean match(uint32 addr, uint32 entry, uint32 mask) { // enter code here }
	Listing 1: Matching function
	d)* Complete the match function of Listing 1.
ļ	

e)* Explain why the routing entry of the default gateway is typically using a subnet mask of 0.	0 1
For the following suproblems, we perform the LPM using a list of routing table entries as a routing table. This list contains a number of n routing table entries. LPM iterates over the list of routing table entries sequentially to	LL1 °
determine the longest matching prefix.	
f)* LPM selects the "longest" entry. What does that mean?	0 1
g) Briefly explain, how many matches must be performed for an unsorted list with n unique entries, before the longest entry can be selected.	0 1
For the following subproblem, we sort the routing table using the IPv4 address values . The list is sorted in ascending order, i.e., the list begins with the routing entry containing the lowest address value. h) Briefly explain, how many matches must be performed for this address-sorted list with n unique entries, before	 10
the longest entry can be selected.	1 2
For the following subproblem, we sort the routing table using the IPv4 subnet masks . The list is sorted in ascending order, i.e., the list begins with the routing entry containing the lowest subnet mask value.	
i) Briefly explain, how many matches must be performed for this mask-sorted list with n unique entries, before the longest entry can be selected.	0 1 2

Problem 3 Load Balancing (15 credits)

Load balancing is important to properly serve large amounts of clients with a good user experience. The following problem is based on the network shown in Figure 3.1.

A company operates two load blancers (LB1 and LB2) each connected to multiple content servers. Furthermore, they operate a nameserver (NS) used to resolve company domains to both load balancers.

On the opposite side, two clients (C1 and C2) frequently contact the company and access their content and a public DNS resolver (R) is available.

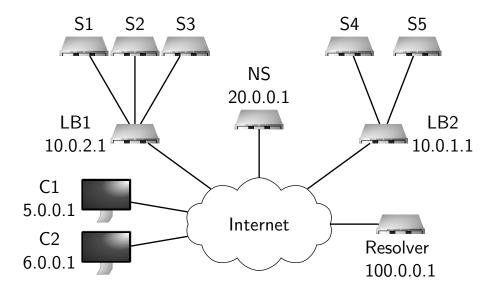


Figure 3.1: Important company with exam relevant content



a)* The company nameserver's zone file contains resource records listed in Table 3.1. Fill in the missing information, so that the nameserver can actively balance the load in between the load balancers.

-	Table 3.1: Company NS Resource Records						
relevant.exam	900	IN					
relevant.exam	900	IN					

	relevant.exam	900	IN			
° H	b)* Explain why it	is not	a god	od idea for the company nameserve	er to return both records in	n one reply.
1 ##						
0	c)* If C1 and C2 u	ıse the	publi	c DNS resolver, why is geographical	load balancing not possil	ole with traditional DNS?
			•			

e) What is the downtime (in minutes) of the service in the worst case if one of the load balancer fails?	2
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	0
f)* Based on the lecture, how can LB1 use modulo hashing to load balance connections to the content servers?	0
g) What happens if one of the content servers fails especially using modulo hashing?	0
h)* How does consistent hashing help with the previous problem?	0 1
i) Assume consistent hashing is used but a planned update needs to be done on a content server. Describe what needs to be done to not disrupt services while using all available servers as long as possible.	0 1 2
j) Instead of DNS load balancing, how else could traffic be load balanced between LB1 and LB2 using routing? Does it prevent service disruptions if a load balancer fails?	0 1 2

Problem 4 Network Calculus (12.5 credits)

This problem investigates performance bounds in networks using Network Calculus.

Consider the following network topology with flow and server descriptions. Flow f_1 traverses Servers s_1 , s_2 , and s_3 . Flow f_2 traverses Servers s_2 and s_3 . Assume each server handles flows according to strict priority scheduling with preemption. Assume Flow f_1 has a low priority and Flow f_2 has a high priority. In the following, we want to apply the Separate Flow Analysis to compute an end-to-end delay bound for Flow f_1 .

Hint: Use the following formula to calculate a left-over service curve: $\beta^{l.o.} = \left[\beta_{R,T} - \gamma_{r,b}\right]^+ = \beta_{R-r,\frac{b+R-T}{R-r}}$. An output arrival curve is given by $\alpha^* = \alpha_{r,b+r\cdot T}$.

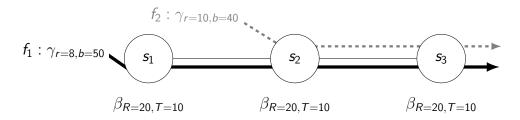
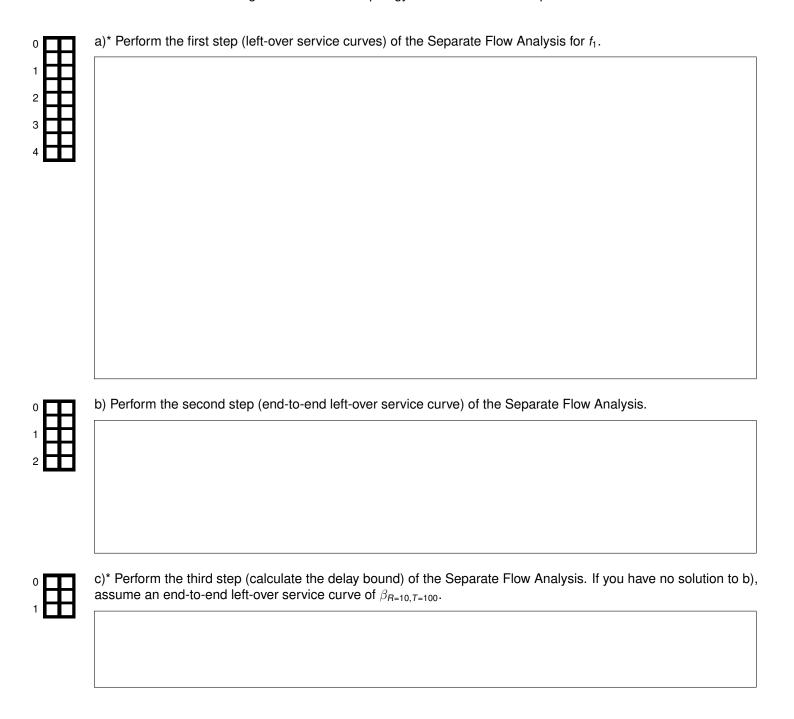


Figure 4.1: Network topology with flow and server specifications



d)* Consider the following scenario to the network in Figure 4.1:	0
• The rate of server s_1 is set to $R = 10$	
Explain the influence on the delay bound of flow f_1 as determined by the Separate Flow Analysis.	ш
e)* Consider a token-bucket constrained flow, traversing a single rate-latency server. The corresponding arrival-and service curves are shown in Figure 4.2. Determine the burst value of the flow and the rate of the server .	0 1 2
Figure 4.2: Arrival- and service curve	
f)* Consider the following statement:	0
• A rate-latency service curve is defined as $\beta_{R,T}(s) = R \cdot s + T$ with rate parameter R and latency parameter T .	
Argue whether or not the statement is correct.	
g)* Name one mathematical framework that can be used to obtain soft real-time guaratees and cannot be used to obtain hard real-time gurantees.	— 0
obtain natu real-time gurantees.	Ш₁

Problem 5 Wireshark (18.5 credits)

The ISO-OSI model defines seven layers in a communication system. For each layer multiple protocols exist. In this problem, a frame is analyzed, referring to the involved protocols. You are given a hexdump of an Ethernet frame including FCS, starting with the Ethernet header. For the following problems we already marked the headers of different layers in Figure 5.1.

					L	ayer	2 He	ader	(Etherne	t)						
0x0000	00	0с	6c	0a	ce	ce	00	25	90	57	22	4a	08	00	45	02
		Layer 3 Header														
0x0010	00	4d	14	6a	40	00	3d	06	b1	18	0a	00	00	0a	14	00
0x0020	00	14	01	bb	d8	dd	04	07	70	b4	50	ba	95	24	80	18
							L	ayer	4 Heade	r						
0x0030	01	f5	d2	90	00	00	01	01	08	0a	3a	с7	db	25	75	df
0x0040	9b	06	17	03	03	00	14	a2	8a	9f	f9	56	2b	a7	54	ac
0x0050	26	ea	07	89	4c	79	0c	96	47	77	СС	af	f2	d1	01	

Figure 5.1: Hexdump of an Ethernet frame including FCS

In this problem you **always** have to substantiate your answers using the bytes in the hexdump in Figure 5.1. You can **either** mark the corresponding bytes directly in the figure **or** list the locations of the corresponding bytes using [...]. Example: the **three** bytes from position 0 to 2 can be written as [0, 2] = 0x000c6c.

0	a)* What is the name of Layer 2 and Layer 3 in the OSI model?
2	Layer 2: Layer 3:
0	b)* In Figure 5.1, mark the FCS.
0	c)* What is the size of the Layer 3 PDU in bytes.
1 2	
0	d)* What is the size of the Layer 4 SDU in bytes.
1 2	

e) Compute the share of the L4 SDU in the whole hexdump.	0
	Ш
f)* Identify the Layer 3 protocol.	0 1
g) List all addresses contained in the Layer 3 header in the common notation.	0 1 2
h) Identify the Layer 4 protocol.	0 1
i) Make an educated guess which protocol can be expected on top of Layer 4.	0 1
j)* Can the two subnets 10.0.0.0/8 and 11.0.0.0/8 be merged? If yes, give the resulting network address and subnet.	0 1
k)* Name and briefly explain two design goals of QUIC, which target weaknesses present with TCP/TLS.	0 1 2

Additional space for solutions-clearly mark the (sub)problem your answers are related to and strike out invalid solutions.

