

#### Note:

- During the attendance check a sticker containing a unique code will be put on this exam.
  This code contains a unique number that associates this exam with your registration
- number.This number is printed both next to the code and to the signature field in the attendance
- check list.

# **Advanced Computer Networking**

Exam:IN2097 / RetakeExaminer:Prof. Dr.-Ing. Georg Carle

**Date:** Wednesday 3<sup>rd</sup> April, 2024 **Time:** 08:00 – 09:15

#### Working instructions

Eexam

Place student sticker here

- This exam consists of **12 pages** with a total of **5 problems**. Please make sure now that you received a complete copy of the exam.
- The total amount of achievable credits in this exam is 75 credits.
- · Detaching pages from the exam is prohibited.
- · Allowed resources:
  - one analog dictionary English  $\leftrightarrow$  native language
- · Subproblems marked by \* can be solved without results of previous subproblems.
- Answers are only accepted if the solution approach is documented. Give a reason for each answer unless explicitly stated otherwise in the respective subproblem.
- · Do not write with red or green colors nor use pencils.
- Physically turn off all electronic devices, put them into your bag and close the bag.

Left room from	to	/ Early submission at
----------------	----	-----------------------

## Problem 1 Quiz (17 credits)

The following questions cover multiple topics and can be solved independently of each other. The multiple choice questions need to be filled out as follows:

Mark correct answers with a cro To undo a cross, completely fill To re-mark an option, use a hur	out the answer option
All multiple choice questions have exactly a single corr	rect answer.
a)* Which of the following is a correct IPv4 address in t	the common notation?
20.255.20.255	20.255.20.256
14.FF.14.FF	20:255:20:255
b)* Which of the following statements about TCP CUB	IC is correct?
It is more robust to packet loss than TCP BBR.	It tries to keep 1 BDP of data inflight.
It is a delay-based congestion control algo- rithm.	It keeps buffers on the path full.
c)* Which of the following statements about adresses i	is correct?
The address space of Ethernet MAC adresses is	<b>bigger</b> than the address space of <b>IPv6</b> addresses.
Ethernet uses a hierarchy for its MAC addresses	
IPv4 uses a hierarchy for its addresses.	
The address space of Ethernet MAC adresses is	s smaller than the address space of IPv4 addresses.
d)* Which of the following statements about Software-I	Defined Networks (SDNs) is correct?
Not all network deployments using Network Fund	ction Virtualization (NFV) are SDNs.
Only software switches can be used in SDNs.	
In an SDN context, all middleboxes are called rou	uter.
The P4 language is the successor of SDN.	
e)* Which of the following statements about DNS zone	es is <b>not</b> correct?
It starts and ends with a SOA record.	
It has its own NS records.	
It must have an A or AAAA records.	
It must have a SOA record.	
f)* How many entries has TBL24 in the DIR-24-8 routir	ng algorithm?
<b>65 536</b>	
Size of TBL24 depends on the number of routing	table entries.
16777216	
4 294 967 296	

h)\* Name one approach from the lecture to detect QUIC deployments besides ZMap.

i)\* NATs can artificially increase the IPv4 address space. Name and shortly explain one fundamental problem due to NAT discussed during the lecture.

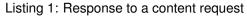
j)\* Name and shortly explain two used IPv6 address types from the lecture.

k)\* What limitation do CNAME records have and why is that relevant for the zone apex?

I)\* Argue what type of load balancing is indicated by the request response shown in Listing 1.

1 HTTP/1.1 301 Moved Permanently

- Location: http://www.google.de/
- 3 Content-Type: text/html; charset=UTF-8
- 5 <htps://dealer.com/seta/ftp-equiv="content-type" content="text/html;charset=utf-8"> <TITLE>301 Moved</TITLE></HEAD><BODY>
- 7 <H1>301 Moved</H1>
- The document has moved
- 9 <A HREF="http://www.google.de/">here</A>. </BODY></HTML>



m)\* Convert the IPv6 address 2001:0db8:0000:0a2e:0000:0370:7334 to its shortened form.



	1

	0
	1
	2

		0
		1
		2

	(	)
	1	I
	1	

		0
		1
		2

### Problem 2 Traceroute and DNS (17 credits)



a)\* What problem does an authoritative name server have, when trying to perform load balancing based on the client's location without any DNS extension information?

0	
1	нн

b) Name and shortly explain the solution for the problem in Subproblem a) presented in the lecture.



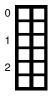
c)\* In order to perform a traceroute to a domain name we need to resolve it to an IP address using DNS. In this problem, we will use IPv4 for our traceroute command. Which record type do we need to query for that?

- 1 ns1.caida.org. 192.172.226.78
  - jungle.caida.org. 192.172.226.32
- *2* jungle.caida.org. 1*3* ns1.ucsd.edu.
  - 192.172.226.32 128.54.16.2

Table 2.1: caida.org NS names and their corresponding IPv4 addresses



d)\* We plan to perform our traceroute towards *caida.org*. The zone of *caida.org* has three NS records. In Table 2.1, we list the names and the corresponding IP addresses. Why are the first two names (*ns1.caida.org* and *jungle.caida.org*) not enough according to the lecture?



e)\* Shortly explain the basic principle of traceroute and why it works.

0		
1		

f)\* Name a reason why the traditional traceroute is not sufficient to always produce consistent results in modern day networks.

1	188.1.37.89	cr-gar1-be2-147.x-win.dfn.de	1ms	1ms	1ms
2	62.40.124.217	dfn.rt1.fra.de.geant.net	9ms	9ms	9ms
3	62.40.98.23	<pre>ae4.mx1.lon.uk.geant.net</pre>	21ms	21ms	21ms
4	163.253.1.118	<pre>core1.ashb.net.internet2.edu</pre>	165ms	163ms	164ms
5	163.253.1.114	<pre>core1.losa.net.internet2.edu</pre>	163ms	163ms	
	137.164.26.200	hpr-lax-agg10-i2.cenic.net	162ms		
6	137.164.26.43	hpr-100ge-sdg-hpr3.cenic.net	162ms	162ms	162ms
7	192.172.226.78	rommie.caida.org	165ms	165ms	165ms

Table 2.2: Simplified traceroute starting from the TUM university network and targeted to caida.org.

g)\* Mark location hints in at least four of the domain names in Table 2.2. Alternatively, you can also reference the hop and write down the location hint in the solutionbox.

F		0
E	Н	1
		2

h)\* Can you always trust the location hints in domain names? Briefly explain your answer!

i)\* What is the reason for the high latency increase between Hops 3 and 4? Explain!

j)\* Explain the specialty visible on Hop 5 and why this poses a problem for the interpretation of traceroute results.

k)\* Name and explain one reason, as covered in the tutorial, why the average 6 are lower than the ones of Hop 4.

	0
-	1
	2

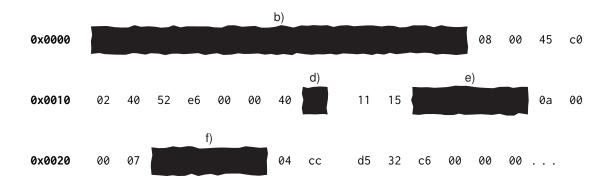
e latencies	of Ho	ps 5	and	6

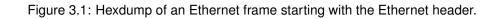
F	L	0	

		0

### Problem 3 Reverse Hexdump (13.5 credits)

This problem investigates a captured Ethernet frame, shown in Figure 3.1. The frame was captured in between  $H_1$  and R as shown in Figure 3.2. For each interface, the MAC and IPv4 address is noted. Unfortunately, the frame was corrupted during the capture process. The lost bytes are marked with black areas in Figure 3.1.





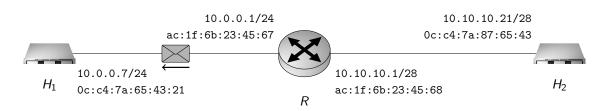


Figure 3.2: Network Topology

As you want to repair the captured hexdump, you have to reconstruct the corrupted parts. If a solution box contains a table, fill each cell with a single byte. You know that the captured frame belongs to an HTTP/3 over QUIC connection between  $H_1$  and  $H_2$ , where  $H_1$  is the client. You also know that the frame went in the direction indicated by an arrow.

0	
1	

a) In our case, only the captured frame was corrupted. If a bit-error occurred during the transmission, how could the receiver possibly detect it? Name a header field that is used for this purpose and explain how the value of this field is determined.

0		
1		
2		

Corrupted block:							

b) Reconstruct the first corrupted block of the captured frame, marked with an b).



c)\* Which **Layer 3** protocol is used in the captured frame? Please also mark and name the field in the hexdump that supports your answer.

#### d) Reconstruct the second corrupted block marked with d) and explain why you chose the value.

Corrupted block:	
Reason:	

e) Reconstruct the third corrupted block marked with e) and name the header field.

				ŀ	_	
Corrupted block:			Header Field:	L		
·						

f) Which field(s) were present in the fourth corrupted block of the captured frame, marked with f)? For each field, state whether you are able to reconstruct the original content with the given information or not. If no, state the reason. If yes, provide the reconstructed content.

g)\* Before  $H_1$  was able to send the first frame to R, it had to resolve the IP address of R to a MAC address. Which protocol was likely used for this purpose?

h) In Subproblem b), you reconstructed the first 12 bytes of the Ethernet header. What would the reconstructed bytes look like in a frame sent from  $H_1$  to R to resolve the MAC address of R for the first time using the protocol from Subproblem g)?

First 12 bytes:						

i)\* How many different IP addresses are contained in the subnet between H<sub>2</sub> and R?

Γ

		0
		1

	Γ	0
		1

## Problem 4 Network Calculus (14 credits)

This problem investigates delay bounds in networks using Network Calculus.

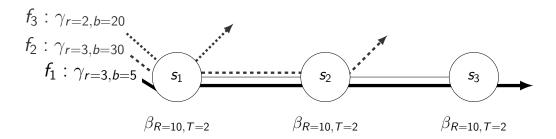
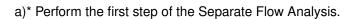
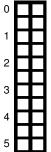


Figure 4.1: Topology with server- and flow specifications

Consider the topology in Figure 4.1. Assume each server employs strict priority queuing. Flow  $f_1$  has the lowest priority while Flows  $f_2$  and  $f_3$  have the highest priority. Flow  $f_1$  traverses three servers, Flow  $f_2$  traverses two, and Flow  $f_3$  traverses one.

We are interested in calculating the end-to-end delay bound of **Flow**  $f_1$  using the Separate Flow Analysis. **Hint:**  $\beta^{l.o.} = \beta_{R-r, \frac{b+R-T}{B-r}}$  and  $\alpha^* = \gamma_{r,b+r,T}$ 







b) Perform the second step of the Separate Flow Analysis.



c) Perform the third step of the Separate Flow Analysis.

d) Assume the following changes to the scenario in Figure 4.1:

- The rate of Flow  $f_1$  is set to r = 4
- The burst of Flow  $f_1$  is set to b = 500

Compute the new delay bound of Flow  $f_1$  under the Separate Flow Analysis.

e)\* A flow with arrival curve  $\gamma_{r=5,b=10}$  is traversing a server with service curve  $\beta_{R=2,T=8}$ . Calculate the **delay** bound of the flow.

f)\* A flow with arrival curve  $\gamma_{r=5,b=480}$  is traversing 1 000 servers connected in series, each with the same service curve  $\beta_{R=20,T=2}$ . Calculate the **delay bound** of the flow. The method you choose should calculate a tight delay bound.

stems g v

g)* Name one approach, bes	sides deterministic	network calculus,	that is suitable for	the analysis of sys
with hard real-time criticalitie	es.			

70 -

	10	
	60	Service curve
	50	
Data [bit]	40	
Data	30	
	20	
	10	
	0	•••
	(	0 10 20 30 40 50 60 70 80 90 100
		Time [s]

Figure 4.2: Arrival- and service curve

h)\* Consider the two curves in Figure 4.2. What is the delay bound of a flow with this non-token-bucket arrival curve traversing a server with this rate-latency service curve?



I			0
			1





		0
	_	1

### Problem 5 Software-Defined Networking (13.5 credits)

This problem investigates a Software-Defined Network (SDN) powered by P4. For the following problems, consider the network given in Figure 5.1. Server 1 is configured to create and only accept tagged frames (VLAN ID 31), Server 2 only creates and accepts untagged frames. PCP and DEI are always set to 0. Switch A is a P4 switch handling the VLAN functionality. In the following subproblems, the P4 program running on Switch A is investigated. Listing 2 shows parts of the used P4 program.

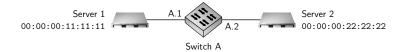


Figure 5.1: Network topology

```
{ bit <48> dstAddr;
    header eth_t
2
                                   bit <48> srcAddr;
                                   bit <16> etherType; }
4
    header veth_ext_t
                                 { bit <3> pcp;
                                   bit <1> dei;
6
                                   bit <12> vid;
                                   bit <16> etherType; }
8
    struct standard_metadata_t { bit <16> ingress_spec;
                                   bit <16> egress_spec;
10
                                 }
                                   / *
12
    struct meta
                                     unused */
                                   eth_t eth;
    struct headers
14
                                   veth_ext_t veth_ext; }
16
    parser ParserImpl(packet_in packet, out headers hdr, inout meta meta, inout standard_metadata_t
        std_meta) {
      // to be defined in Subproblem a)
18
    }
20
    control Pipeline (inout headers hdr, inout metadata meta, inout standard_metadata_t std_meta) {
      action drop() {
22
        mark_to_drop();
      }
24
      action untag(bit <16> egress) {
        std_meta.egress_spec = egress;
26
        hdr.eth.etherType = hdr.veth_ext.etherType;
        hdr.veth_ext.[...]; // to be defined in Subproblem b)
28
      action tag(bit <16> egress, bit <12> vid) {
30
        std_meta.egress_spec = egress;
        hdr.veth_ext.etherType = hdr.eth.etherType;
32
        hdr.eth.etherType = 0x8100;
        hdr.veth_ext.[...]; // to be defined in Subproblem c)
34
        hdr.veth_ext.pcp = // to be defined in Subproblem c)
        hdr.veth_ext.dei = // to be defined in Subproblem c)
        hdr.veth_ext.vid = // to be defined in Subproblem c)
36
      table forward {
38
        actions = {
40
          tag:
          untag:
42
          drop;
44
        key = {
          std_meta.ingress_spec: exact;
46
        ļ
        size = 4:
48
        default_action = drop;
      }
50
      apply {
        if (hdr.eth.isValid()) {
52
          forward.apply();
        }
54
      }
    }
56
    \\ ...
```

Listing 2: VLAN P4 program

a)\* Write a parser using the P4 language for the P4 program in Listing 2. The parser starts at the start state and must be able to accept tagged and untagged Ethernet frames.

b)\* Complete the untag() action of Listing 2 (Line 27).

 $\verb+hdr.veth\_ext.$ 

c)\* Complete the tag() action of Listing 2 (Lines 33–36).

hdr.veth\_ext.

hdr.veth\_ext.pcp =

hdr.veth\_ext.dei =

hdr.veth\_ext.vid =

d)\* The P4 program cannot work correctly without table data containing correct forwarding rules. Give the rules for Switch A, to correctly encapsulate and forward frames of Servers 1 and 2. Use the information given in Figure 5.1.

Key	Action	Action data	
-			

		0
		1
		2
		3
L		
┝	Н	4

		0
		1
	L	
		2

	0
E	1
	2
	3
	4
	5

Additional space for solutions-clearly mark the (sub)problem your answers are related to and strike out invalid solutions.

